06-01-00

A

JC830 U.S. PTO 05/31/00

JC658 U.S. PTO 09/584516 05/31/00

UTILITY PATENT APPLICATION TRANSMITTAL

Submit an original and a duplicate for fee processing

(Only for new nonprovisional applications under 37 CFR 1.53(b))

ADDRESS TO: **Assistant Commissioner for Patents** **Box Patent Application** **Washington, D.C. 20231**	Attorney Docket No. 00,011 First Named Inventor David A. Grabelsky Express Mail No. EL007258559US Total Pages (including transmittal)

APPLICATION ELEMENTS	ACCOMPANYING APPLICATION PARTS
1. ☒ Transmittal Form with Fee 2. ☒ Specification (including claims and abstract) [Total Pages 40] 3. ☒ Drawings [Total Sheets 12] 4. ☐ Oath or Declaration [Total Pages] a. ☐ Newly executed b. ☐ Copy from prior application **[Note Boxes 5 and 17 below]** i. ☐ <u>Deletion of Inventor(s)</u> Signed statement attached deleting inventor(s) named in the prior application 5. ☐ Incorporation by Reference: The entire disclosure of the prior application, from which a copy of the oath or declaration is supplied under Box 4b, is considered as being part of the disclosure of the accompanying application and is hereby incorporated by reference therein. 6. ☐ Microfiche Computer Program 7. ☐ Nucleotide and/or Amino Acid Sequence Submission a. ☐ Computer Readable Copy b. ☐ Paper Copy c. ☐ Statement verifying above copies	8. ☐ Assignment Papers 9. ☐ Power of Attorney 10. ☐ English Translation Document (if applicable) 11. ☐ Information Disclosure Statement (IDS) ☐ PTO-1449 Form ☐ Copies of IDS Citations 12. ☐ Preliminary Amendment 13. ☒ Return Receipt Postcard (Should be specifically itemized) 14. ☐ Small Entity Statement(s) ☐ Enclosed ☐ Statement filed in prior application; status still proper and desired 15. ☐ Certified Copy of Priority Document(s) 16. ☐ Other:

17. **If a CONTINUING APPLICATION**, check appropriate box and supply the requisite information:
☐ Continuation ☐ Divisional ☒ Continuation-in-part of prior application Serial No. 09/035,600.

09584516 .053100

APPLICATION FEES				
BASIC FEE				$690.00
CLAIMS	NUMBER FILED	NUMBER EXTRA	RATE	
Total Claims	- 20 =		x $18.00	$
Independent Claims	- 3 =		x $78.00	$
☐ Multiple Dependent Claims(s) if applicable			+$270.00	$
			Total of above calculations =	$
			Reduction by 50% for filing by small entity =	$( )
☐ Assignment fee if applicable			+ $40.00	
			TOTAL =	$ 690.00

18. ☐ Please charge my Deposit Account No. 13-2490 in the amount of $.

19. ☒ A check in the amount of $ .00 is enclosed.

20. The Commissioner is hereby authorized to credit overpayments or charge any additional fees of the following types to Deposit Account No. 13-2490:
    a. ☒ Fees required under 37 CFR 1.16.
    b. ☒ Fees required under 37 CFR 1.17.
    c. ☒ Fees required under 37 CFR 1.18.

21. ☐ The Commissioner is hereby generally authorized under 37 CFR 1.136(a)(3) to treat any future reply in this or any related application filed pursuant to 37 CFR 1.53 requiring an extension of time as incorporating a request therefor, and the Commissioner is hereby specifically authorized to charge Deposit Account No. 13-2490 for any fee that may be due in connection with such a request for an extension of time.

22. CORRESPONDENCE ADDRESS	
Name	**McDonnell Boehnen Hulbert & Berghoff**
Address	32nd Floor, 300 South Wacker Drive
City, State, Zip	Chicago, Illinois 60606

23. SIGNATURE OF APPLICANT, ATTORNEY, OR AGENT REQUIRED	
Name	S. Richard Carden Reg. No. 44,588
Signature	
Date	May 31, 2000

UTILITY (Rev. 11/18/97)

CERTIFICATE OF MAILING

PATENT

Express Mail No. EL007258559US

Deposited: May 31, 2000

I hereby certify that the attached correspondence, identified below, is being deposited with the United States Postal Service as "Express Mail Post Office to Addressee" under 37 CFR 1.10 on the date indicated above and is addressed to, Box Patent Application, Assistant Commissioner for Patents, Washington, D.C. 20231.

By: Wendy O'Neal

In re the Application of: David A. Grabelsky, Michael S. Borella, John Poplett and Richard J. Dynarski.

Docket No. MBHB00-011

For: METHOD FOR ADDRESS MAPPING IN A NETWORK ACCESS SYSTEM AND A NETWORK ACCESS DEVICE FOR USE THEREWITH

X Transmittal Letter

X Specification (40 sheets including claims and abstract)

X Drawings (12 sheets)

____ Declaration and Power of Attorney

X Filing Fee (Large Entity)

X Return Post Card

APPLICATION FOR UNITED STATES LETTERS PATENT
IN THE UNITED STATES PATENT AND TRADEMARK OFFICE
(3Com Case No. 2876.CS.US.P; MBHB Case No. 00,011)

Title: METHOD FOR ADDRESS MAPPING IN A NETWORK ACCESS SYSTEM AND A NETWORK ACCESS DEVICE FOR USE THEREWITH

Inventors: David A. Grabelsky
3800 Lee Street
Skokie, Illinois 60076

Citizen of the United States of America

Michael S. Borella
1208 Haverhill Circle
Naperville, Illinois 60563

Citizen of the United States of America

John Poplett
1011 Keystone Avenue
River Forest, Illinois 60305

Citizen of the United States of America

Richard J. Dynarski
3 South Sequoia Drive
Glen Ellyn, Illinois 60137

Citizen of the United States of America

Assignee: 3Com Corporation
5400 Bayfront Plaza
Santa Clara, California 95052

Attorneys: S. Richard Carden
Bradley J. Hulbert
McDonnell Boehnen Hulbert & Berghoff
300 S. Wacker
32nd Floor
Chicago, IL 60606

METHOD FOR ADDRESS MAPPING IN A NETWORK ACCESS SYSTEM AND A NETWORK ACCESS DEVICE FOR USE THEREWITH

RELATED APPLICATIONS

This application is a continuation in part of U.S. Application Serial No. 09/035,600, filed March 15, 1998.

BACKGROUND OF THE INVENTION

A. Field of the Invention

The present invention relates to computer networks. More specifically, the invention relates to a method for assigning a public Internet Protocol ("IP") address from an address server on a network to a Realm Specific Internet Protocol aware host on the network having a private IP address.

B. Description of the Related Art

The IP is an addressing protocol designed to route traffic within a network or between networks. Current versions of IP such as IP version 4 ("Ipv4") are becoming obsolete because of limited address space. With a 32-bit address-field, it is possible to assign 2^{32} different addresses, which is 4,294,967,296, or greater than 4 billion possible addresses. A unique IP number is typically assigned to network devices on a network using IP, whether or not the network is connected to the Internet. Most organizations, such as corporations and universities have multiple networks using IP, with multiple network devices assigned an IP address. With the explosive growth of the Internet and intranets, IP addresses using a 32-bit address-field may soon be exhausted. IP version 6 ("IPv6") proposes the use of a 128-bit address-field for IP addresses.

McDonnell Boehnen Hulbert & Berghoff
300 South Wacker Drive, 32nd Floor
Chicago, IL 60606
(312) 913-0001

However, a large number of networks including a large number of Internet nodes will still be using older versions for IP with a 32-bit address space for many years to come.

The sharing of a public IP address among multiple hosts is a useful method in cases where a local area network (LAN) is comprised of multiple IP hosts, but possess only a limited number of public IP addresses that reside at a router. Such a network is referred to as a stub network. Each local host on the stub network has only a private (internal) IP address. When communicating amongst themselves, local hosts use their local IP addresses. However, for communications with the public (external) IP network, some form of address mapping/sharing must be implemented in order to allow the local hosts to send/receive packets to/from entities on the external IP network.

Network access systems ("NAS") generally consist of multiple device subsystems, such as modem cards, which reside in a chassis, and are connected by one or more internal communications systems, such as a communications bus. A larger NAS may consist of multiple chassis connected in a LAN, or some local communications system. Typically, the entire NAS provides one or a few public IP interfaces, usually associated with a router card or subsystem, for communications with the external IP network. Devices on the external IP network can communicate with the NAS through these one or a few interfaces. In certain cases, it is desirable to partition and group the internal NAS resources in such a way to make it appear as multiple, virtual NAS systems to devices on the external IP network. It may even be desirable to make each internal device subsystem individually addressable on the external, public IP network. One way to accomplish this is to implement an IP stack on each internal device subsystem, connect them with an internal LAN, and provide external access by incorporating routing functionality at the NAS's external IP interface. In such a configuration, the internal IP network may also

McDonnell Boehnen Hulbert & Berghoff
300 South Wacker Drive, 32nd Floor
Chicago, IL 60606
(312) 913-0001

provide the internal communications for the system, or augment some other bus-like system. If the IP addresses of the internal device subsystems are only private, then the internal LAN can be viewed as a stub network. In this case, some form of address mapping/sharing must be implemented to enable the internal, subsystem devices to communicate with the external IP network as IP devices.

Network address translation ("NAT") has been proposed to extend the lifetime of Internet Protocol ("IP") version 4 ("IPv4") and earlier versions of IP by allowing a network to exist behind a set of public IP addresses. *See* P. Srisureh, "IP Network Address Translator (NAT) Terminology and Considerations," IETF RFC 2663, Aug. 1999, which is incorporated herein by reference. NAT provides a method for transparent bi-directional communication between a private routing realm, for example a private intranet, and an external routing realm, for example, the Internet. Through use of NAT, addresses of packets sent by the first realm are translated into addresses associated with the second realm. Use of private IP addresses in conjunction with a NAT implementation in a network address server allows the ISP to conserve globally-routable public IP addresses. When a device or node using private addressing desires to communicate with the external world, a private address is translated to a common public IP address used for communication with an external network by a NAT device.

There are several problems associated with using NAT to extend the life of IP. NAT interferes with the end-to-end routing principle of the Internet that recommends that packets flow end-to-end between network devices without changing the contents of any packet along a transmission route. (see e.g., Routing in the Internet, by C. Huitema, Prentice Hall, 1995) Problems with NAT support for end-to-end protocols, especially those that authenticate or encrypt portions of data packets, are particularly well-known. *See, e.g.*, Holdrege et al.,

McDonnell Boehnen Hulbert & Berghoff
300 South Wacker Drive, 32nd Floor
Chicago, IL 60606
(312) 913-0001

"Protocol Complications with the IP Network Address Translator," Internet Draft <draft-ietf-nat-protocol-complications-01>, Jun. 1999. In applications that transmit IP addresses in packet payloads, NAT requires an application layer gateway to function properly. NAT also creates difficulties when applied to Internet security applications.

Current versions of NAT replace a private network address in a data packet header with an external network address on outbound traffic, and replace an external address in a data packet header with a private network address on inbound traffic. This type of address translation is computationally expensive, causes security problems by preventing certain types of encryption from being used, or breaks a number of existing applications in a network that cannot do NAT (e.g., File Transfer Protocol ("FTP")).

Current versions of NAT also may not gracefully scale beyond a small network containing a few dozen nodes or devices because of the computational and other resources required. NAT potentially requires support for many different internal network protocols be specifically programmed into a translation mechanism for external protocols in a NAT device such as a NAT router. As is known in the art, a router translates differences between network protocols and routes data packets to an appropriate network node or network device. Computational burdens placed on a NAT router may be significant and degrade network performance, especially if several NAT-enabled stub networks share the same NAT router. In a worst case scenario, a NAT router translates every inbound and outbound data packet.

Realm Specific Internet Protocol (RSIP) has been proposed as an alternative for NAT. *See* M. Borella et al., "Realm Specific IP: Protocol Specification," Internet Draft <draft-ietf-nat-rsip-protocol-06>, Mar. 2000 (hereinafter "RSIP-PROTOCOL"), which is incorporated herein by reference. Using RSIP, a host and a gateway negotiate the use of a public IP address and possibly

some number of Transmission Control Protocol (TCP)/User Datagram Protocol (UDP) ports. As is known in the art, Transmission Control Protocol ("TCP") and User Datagram Protocol ("UDP") are often used over IP in computer networks. TCP provides a connection-oriented, end-to-end reliable protocol designed to fit into a layered hierarchy of protocols that support multi-network applications. UDP provides a transaction oriented datagram protocol, where delivery and duplicate packet protection are not guaranteed. After enabling RSIP in the host, the RSIP-aware host can utilize one or more public IP addresses residing on the RSIP gateway. An important capability of RSIP compared with other methods such as NAT, is that local host devices may terminate IPsec with external IP entities, even while they share the single public address of the LAN router device.

RSIP requires that an application on the host communicate with an application on the RSIP gateway, so the communication link must be configured for IP before this communication occurs. The RSIP client must also know the IP address of the RSIP gateway, so that it can contact the gateway directly. Thus, there is a need in the art for a method by which an RSIP host can determine the IP address of an RSIP gaetway. There is also a need in the art for a process by which a network device and a network server may use RSIP to assign public IP addresses from the server to the network device in order to conserve globally-routable IP addresses.

McDonnell Boehnen Hulbert & Berghoff
300 South Wacker Drive, 32nd Floor
Chicago, IL 60606
(312) 913-0001

SUMMARY OF THE INVENTION

In accordance with an illustrative embodiment of the present invention, the problems associated with NAT are overcome. A device and method for implementing Realm Specific Internet Protocol ("RSIP") in a network access system is provided.

In accordance with a first aspect of the invention, a network access address mapping system is provided. The network access address mapping system includes a plurality of first network subdevices, and a second network subdevice. The plurality of first network subdevices is connected on a first network by a common communications path. Each of the first network subdevices has a private network address for communicating with the network subdevices on the first network. The second network subdevice has a private network address, a public network address, and one or more ports. The private network address of the second network subdevice is used for communication with the network subdevices on the first network. The public network address of the second network subdevice is used for communicating with network devices on an external public network. The combination network address includes the public network address and is used for identifying any of the first network subdevices during communication with network devices on the external network. The first network subdevices request allocation of the public network address and one or more ports from the second network subdevice for communication with network devices on the external network.

In one embodiment, the first and second network subdevices and the first network comprise a stub network. In a further preferred embodiment, the first and second network subdevices and the first network comprise a single, self-contained network device. In a further preferred embodiment, the first and second network subdevices comprise cards in a rack having a

McDonnell Boehnen Hulbert & Berghoff
300 South Wacker Drive, 32nd Floor
Chicago, IL 60606
(312) 913-0001

common backplane. In a further preferred embodiments, the first network devices in the network access server are positioned in a plurality of chassis.

In an exemplary embodiment of the invention, the first and second network devices are Internet protocol-addressable devices, the private and public addresses are Internet protocol ("IP") addresses, and the first and second networks are IP networks. In this embodiment, the public IP address is a globally unique IP address.

In accordance with a second aspect of the invention, a method for implementing RSIP in the network access system of the invention is provided. The method of the invention comprises the steps of requesting by a first network subdevice having a private network address a public network address and one or more ports from a second network subdevice having a private network address and a public network address; receiving the public network address and the one or more ports by the first network subdevice from the second network subdevice; updating an address-to-address table maintained in the second network subdevice to reflect allocation of the public network address and the one or more ports to the first network subdevice, and creating a combination network address comprising the public network address and the one or more ports to identify the first network subdevice during communications with network devices on an external network.

The foregoing and other aspects and advantages of illustrative embodiments of the present invention will be more readily apparent from the following detailed description, which proceeds with references to the accompanying drawings.

McDonnell Boehnen Hulbert & Berghoff
300 South Wacker Drive, 32^{nd} Floor
Chicago, IL 60606
(312) 913-0001

BRIEF DESCRIPTION OF THE DRAWINGS

Presently preferred embodiments of the invention are described with reference to the following drawings, wherein:

FIG. 1 is a block diagram illustrating a network access system using Realm Specific Internet Protocol;

FIG. 2 is a block diagram illustrating a preferred embodiment of the network access system of FIG. 1;

FIG. 3 is a block diagram illustrating a protocol stack for a network subdevice;

FIG. 4 is a block diagram illustrating a group of Realm Specific Internet Protocol messages;

FIG. 5 is a block diagram illustrating a Realm Specific Internet Protocol message layout;

FIG. 6 is a block diagram illustrating a register request message layout;

FIG. 7 is a block diagram illustrating a register response message layout;

FIG. 8 is a block diagram illustrating an assign request message layout;

FIG. 9 is a block diagram illustrating an assign request message layout for an embodiment in which the RSIP type is RSA-IP;

FIG. 10 is a block diagram illustrating an assign request message layout for an embodiment in which the RSIP type is RSAP-IP;

FIG. 11 is a block diagram illustrating an assign response message layout;

FIG. 12 is a block diagram illustrating an assign response message layout for an embodiment in which the RSIP type is RSA-IP;

FIG. 13 is a block diagram illustrating an assign response message layout for an embodiment in which the RSIP type is RSAP-IP;

FIG. 14 is a block diagram illustrating a combination network address;

FIG. 15 is a block diagram illustrating a port-to-internal address table;

FIG. 16 is a flow diagram illustrating a method for creating a combination network address; and

FIG. 17 is a flow diagram illustrating a method for implementing RSIP in the network access server.

DETAILED DESCRIPTION OF THE PRESENTLY PREFERRED EMBODIMENTS

Figure 1 shows the high-level architecture of a network access system ("NAS") **2**. In the figure, NAS **2** comprises a number of individually addressable first network subdevices **6** having a common communications path **8** connected together by an internal private network **10**. External public addressability is provided via one or a few public interfaces on a second network subdevice **7**. Network subdevices **6** and **7** communicate with an external public network **14**, such as the Internet, or a second private network, via the public interfaces **12** using one or a few a common globally unique network address **44**. The invention, however, is not limited to these external networks, and those of skill in the art will recognize the utility of the invention for transmission of data over any packet based network. Communications between the internal private network **10** and the external public network **14** may take place over the public-switched telephone network ("PSTN"), a cable television network, or any other suitable network or medium.

In a preferred embodiment, NAS **2** comprises a first network subdevice **6** having a private network address **9** that requests the assignment of a public network address **11** from a second network subdevice **7** using a first protocol **13**. In a preferred embodiment, the internal and external networks **10** and **14** are IP networks, the network subdevices **6** and **7** are IP addressable, the public interfaces **12** are IP interfaces, private and public network addresses **9** and **11** are IP addresses, and the first protocol **13** used for requesting assignment of a public network address **11** is Realm Specific Internet Protocol.

In one preferred embodiment, as shown in Figure 2, the first network subdevice **6** comprises a chassis **18** housing one or a plurality of communications cards **24**, and the second network subdevice **7** comprises a router subsystem **20**. The first network subdevice **6** and the

McDonnell Boehnen Hulbert & Berghoff
300 South Wacker Drive, 32nd Floor
Chicago, IL 60606
(312) 913-0001

second network subdevice **7** communicate via an internal IP network **10**. Each communications card **24** further preferably comprises an IP interface **26**, an RSIP host **28**, a device control application ("DCA") **30**, and a data application **32**. The IP interface **26** on the communication card **24** is connected to the internal IP network **10**. The router subsystem **20** preferably comprises one or more IP interfaces **12**. In one preferred embodiment, the router subsystem **20** comprises three IP interfaces **12**: one connected to the internal IP network **10**, and one each to an external data network **27** and an external IP signaling network **29**. Preferably, NAS **2** is a self-contained unit; however other configurations are possible, and the invention is not limited to this embodiment.

Preferred embodiments of NAS **2** may include additional sub-components and elements, as well as additional public IP interfaces **12**. In particular, reference to any specific network access elements are completely omitted from Figure 2. It should also be understood that the external networks shown in the figure need not be separate networks; nor is there any implied limitation on the number of external networks.

As used herein, the term "RSIP client" refers to an application that runs the client side of the RSIP protocol. As used herein, the term "RSIP host" refers to the physical device where the RSIP client application resides. Referring to Figure 2, the RSIP host of NAS **2** corresponds to a communications card 24. As used herein, the term "RSIP server" refers to an application that runs the server side of the RSIP protocol. As used herein, the term "RSIP gateway" refers to the physical device where the RSIP server application resides. Referring to Figure 2, the RSIP gateway of NAS **2** corresponds to the router subsystem **20**.

An operating environment for network devices and routers of the present invention includes a processing system with at least one high speed Central Processing Unit ("CPU") and a

McDonnell Boehnen Hulbert & Berghoff
300 South Wacker Drive, 32nd Floor
Chicago, IL 60606
(312) 913-0001

memory. In accordance with the practices of persons skilled in the art of computer programming, the present invention is described below with reference to acts and symbolic representations of operations that are performed by the processing system, unless indicated otherwise. Such acts and operations may be referred to as being "computer-executed" or "CPU executed."

It will be appreciated that acts and symbolically represented operations include the manipulation of electrical signals by the CPU. The memory locations where data bits are maintained are physical locations that have particular electrical, magnetic, optical, or organic properties corresponding to the data bits.

The data bits may also be maintained on a computer readable medium including magnetic disks, optical disks, and any other volatile (e.g., Random Access Memory ("RAM")) or non-volatile (e.g., Read-Only Memory ("ROM")) mass storage system readable by the CPU. The computer readable medium includes cooperating or interconnected computer readable medium, which may exist exclusively on the processing system or may be distributed among multiple interconnected processing systems that may be local or remote to the processing system.

In network address translation schemes known in the prior art, the router subsystem **20** translates an internal network address, such as an internal IP address used on the internal IP network **10**, to an external network address such as an IP address for outgoing traffic to the external IP data network **14**. The router subsystem **20** also translates an external network address to an internal network address for incoming traffic from external IP data network **14**. A NAT router assumes the entire computational burden for network address translation. For large stub networks having 50 or more network devices or subdevices, the NAT router may become a

McDonnell Boehnen Hulbert & Berghoff
300 South Wacker Drive, 32nd Floor
Chicago, IL 60606
(312) 913-0001

bottleneck. In the worst case, every packet passing through the NAT router requires address translation.

In an illustrative embodiment of the present invention, Realm Specific Internet Protocol ("RSIP") is used to overcome the difficulties associated with NAT. In a preferred embodiment of the invention, the first network subdevices **6** on the internal IP network **10** request a globally unique public network address from the second network subdevice **7**, as well as a set of locally unique ports, for external communication with the external network **14**, such as an external IP data network. The second network subdevice **7** then creates a combination network address **112** comprising the globally unique network address and the locally unique ports, to identify information transmitted to and from the first network subdevice **6**.

RSIP Protocol Stack

FIG. 3 is a block diagram illustrating a layered protocol stack **46** for a communications card **24** on internal IP network **10** used for RSIP. The layered protocol stack **46** is described with respect to Internet Protocol suites comprising from lowest-to-highest, a link layer **48**, a network layer **50**, a transport layer **52** and an application layer **54**. However, more or fewer layers could alternatively be used, and different layer designations could also be used for the layers in the protocol stack **46** (e.g., layering based on the Open Systems Interconnection ("OSI") model).

The network layer **50** includes an IP layer **58** (hereinafter "IP **58**"), an Internet Group Management Protocol ("IGMP") layer **62**, and a Control Message Protocol ("ICMP") layer **64**, and may also include a RSIP layer **60** (not shown). As is known in the art, the IP **58** is an addressing protocol designed to route traffic within a network or between networks. The IP **58** is

McDonnell Boehnen Hulbert & Berghoff
300 South Wacker Drive, 32nd Floor
Chicago, IL 60606
(312) 913-0001

described in RFC-791, J. Postel, *Internet Protocol*, Sep. 1, 1981, incorporated herein by reference.

Above the network layer **50** is a transport layer **52**. The transport layer **52** includes a Transmission Control Protocol ("TCP") layer **58** and a User Datagram Protocol ("UDP") layer **68**.

The RSIP gateway **38** allocates locally unique ports to a first network subdevice **6** having the RSIP layer **60**. In one embodiment of the present invention, the RSIP layer **60** is a separate protocol layer in the network layer **50**. In another embodiment of the present invention, the RSIP layer **60** is implemented as part of the ICMP layer **64** and is not a separate protocol layer. In yet another embodiment of the present invention, the RSIP layer **60** is run over either a Transmission Control Protocol or User Datagram Protocol. The RSIP layer **60** is explained below.

The IGMP layer **62**, hereinafter IGMP **62**, is responsible for multicasting. For more information on the IGMP **62**, see RFC-2236, W. Fenner, *Internet Group Management Protocol, Version 2*, November 1997, incorporated herein by reference.

The ICMP layer **64**, hereinafter ICMP **64**, is used for Internet Protocol control. The main functions of the ICMP **64** include error reporting, reachability testing (e.g., "pinging"), route-change notification, performance, subnet addressing and other maintenance. For more information on the ICMP **64**, see RFC-950, J.C. Mogul and J. Postel, *Internet Standard Subnetting Procedure*, Aug. 1, 1985, incorporated herein by reference.

The TCP layer **66**, hereinafter TCP **66**, provides a connection-oriented, end-to-end reliable protocol designed to fit into a layered hierarchy of protocols which support multi-network applications. The TCP **66** provides for reliable inter-process communication between pairs of processes in network devices attached to distinct but interconnected networks. For more

information on the TCP **66**, see RFC-793, J. Postel, *Transmission Control Protocol*, Sep. 1, 1981, incorporated herein by reference.

The UDP layer **68**, hereinafter UDP **68**, provides a connectionless mode of communications with datagrams in an interconnected set of computer networks. The UDP **68** provides a transaction oriented datagram protocol, where delivery and duplicate packet protection are not guaranteed. For more information on the UDP **68,** see RFC-768, J. Postel, *User Datagram Protocol*, Aug. 28, 1980, incorporated herein by reference. Both the TCP **66** and the UDP **68** are not required in protocol stack **42**; either the TCP **66** or the UDP **68** can be used without the other.

Above the transport layer **52** is an application layer **54** where application programs to carry out desired functionality for a network device reside.

More or fewer protocol layers may alternatively be used in the protocol stack **42**.

RSIP Protocol

FIG. 5 is a block diagram illustrating the general layout for an RSIP protocol message. The RSIP protocol **70** includes a register request message **72**, a register response message **74**, an assign request message **76** and an assign response message **78**. Each RSIP protocol **70** message is in type-length-value ("TLV") format. Additional messages may also be used for RSIP protocol messages, as described in RSIP-PROTOCOL.

As shown in FIG. 5, each RSIP protocol **70** message comprises a header **80** including three mandatory fields followed by required parameters **82**. The mandatory fields comprise a version field **83**, a message type field **84**, and an overall length field **85**. The version field **83** is one byte and indicates the version of RSIP being used. The message type field **84** is one byte and indicates the specific message type, e.g. a value of 2 may indicate a register request message

McDonnell Boehnen Hulbert & Berghoff
300 South Wacker Drive, 32nd Floor
Chicago, IL 60606
(312) 913-0001

72, 3 may indicate a register response message **74**, etc. The overall length field **85** is two bytes and indicates the length of the entire message, including the header **80**. The required parameters **82** each comprise a one byte code field **86**, a two byte length field **87**, and a variable length value field **88**. The length field **87** specifies the length of the value field **88** only.

In an illustrative embodiment of the present invention, the register request message **72** is sent from the RSIP host **28** to the RSIP gateway **38** to request a globally unique IP address **44** and a block of locally unique ports **42**. FIG. 6 is a block diagram illustrating a register request message **72** layout, which comprises header **80**.

In one embodiment of the present invention, the RSIP host **28** transmits the register request message **72** upon boot. The RSIP protocol **70** can exist as a separate protocol or can be integrated into an known configuration protocol, for example Dynamic Host Configuration Protocol ("DHCP"). DHCP **89** is a protocol for passing configuration information such as IP addresses to hosts on a network. For more information on DHCP **89** see RFC-2131, R. Droms, *Dynamic Host Configuration Protocol*, March 1997, incorporated herein by reference. The format of DHCP **89** messages is based on the format of BOOTP messages described in RFC-951 and RFC-1542, incorporated herein by reference. From a network device's point of view, DHCP **89** is an extension of the BOOTP mechanism.

In another embodiment of the present invention, the RSIP host **28** requests a globally unique IP address **44** and locally unique ports **42** after boot when a protocol layer in layered protocol stack **46** makes an initial request for an external network **14**. The RSIP host **28** may also request a globally unique address **44** and locally unique ports **42** when the number of globally unique addresses **44** and locally unique ports **42** required falls below the number of globally unique addresses and ports allocated.

The register response message **74** is sent from the RSIP gateway **38** back to the RSIP host **28** either confirming or denying the register request message **72**. FIG. 7 is a block diagram illustrating a register response message **74** layout. The register response message **74** comprises header **80** followed by required parameters client ID **90** and flow policy **92**. Register response message **74** may also include optional parameters for RSIP type **94** and tunnel type **96**. Client ID parameter **90** has a code field **86** value of 4 and a length of four bytes.

Flow policy parameter **92** has a code field **86** value of 9 and a length of two bytes, the first byte of which specifies the local flow policy and the second byte of which specifies the remote flow policy. Flow policies are described in greater detail in RSIP-PROTOCOL.

Optional RSIP type parameter **94** has a code field **86** value of 7 and a length of one byte. This byte specifies whether Realm Specific Address IP (RSA-IP) or Realm Specific Address and Port IP (RSAP-IP) will be used. In RSA-IP, the RSIP gateway **38** allocates each RSIP host **28** a globally unique public IP address **44**, and may allocate a number of locally unique ports **42** not associated with the unique public IP address. In RSAP-IP, the RSIP gateway **38** allocates each RSIP host **28** a globally unique public IP address **44** and a number of locally unique ports **44** associated with the address.

Optional tunnel type parameter **96** has a code field **86** value of 6 and a length of one byte. Possible tunnel types specified by this parameter include IP-IP (value of 1), GRE (value of 2), and L2TP (value of 3).

Upon receiving a successful register response message **74**, the RSIP host **28** sends an assign request message **76** to RSIP gateway **38**. FIG. 8 is a block diagram illustrating a assign request message **76** layout. The assign request message **76** comprises a header **80** followed by required parameters client ID **90** and an address and port parameters **98**. As shown in FIG. 9, if

McDonnell Boehnen Hulbert & Berghoff
300 South Wacker Drive, 32nd Floor
Chicago, IL 60606
(312) 913-0001

the system is using RSA-IP, the address and port parameters **98** comprise mandatory parameters for local address **100**, remote address **102**, and remote ports **43**, and optional parameters for lease time **106** and tunnel type **96**. Optional parameters are indicated by brackets. As shown in FIG. 10, if the system is using RSAP-IP, the address and port parameters **98** comprise mandatory parameters for local address **100**, local ports **42**, remote address **102**, and remote ports **43**, and optional parameters for lease time **106** and tunnel type **96**. Optional parameters are indicated by brackets.

The address parameters **100** and **102** have a one byte code field **86** with a value of 1 and a two byte length field **87** that specifies the remaining length of the message. The first byte of the length is a type field and the remaining length is a value field. The length of the value field **88** depends on the type of address selected.

The port parameters **42** and **43** have a one byte code field **86** with a value of 2 and a two byte length field **87** that specifies the remaining length of the message. The first byte of the length is a one byte number field that specifies the number of ports. The remaining length consists of one or more two byte port fields that specify the ports to be allocated.

The lease time parameter **106** has a one byte code field **86** with a value of 3 and a four byte length field **87** that specifies the remaining length of the message. The value in the remaining length specifies the amount of time that the binding will remain active.

The Assign response message **78** is sent from the RSIP gateway **38** back to the RSIP host **28** with a globally unique public IP address and one or more locally unique ports for use by the RSIP host **28**. FIG. 11 is a block diagram illustrating an assign response message **78** layout. The assign response message **78** comprises header **80** followed by required parameters client ID **90**, bind ID **110**, and address and port parameters **98**. As shown in FIG. 12, if the system is using

McDonnell Boehnen Hulbert & Berghoff
300 South Wacker Drive, 32nd Floor
Chicago, IL 60606
(312) 913-0001

RSA-IP, the address and port parameters **98** comprise mandatory parameters for local address **100**, remote address **102**, remote ports 43, lease time **106** and tunnel type **96**. As shown in FIG. 13, if the system is using RSAP-IP, the address and port parameters **98** comprise mandatory parameters for local address **100**, local ports **42**, remote address **102**, remote ports **43**, lease time **106** and tunnel type **96**. Note that the lease time and tunnel type parameters are mandatory in the Assign response message **78**, while they are optional in the assign request message **76**.

Once the RSIP gateway **38** assigns a globally unique public IP address **44** and one or more locally unique ports **42** to the RSIP host **28**, the RSIP host **28** saves the block of locally unique ports **42** that it may use. The one or more locally unique ports **42** are allocated to protocols and applications in layered protocol stack **46** on the RSIP host **28** to replace local or default ports. If no addresses are available, the RSIP gateway **38** returns an error message to the RSIP host **28**. The locally unique ports **42** are saved in a data structure with a flag-field indicating whether the locally unique port **42** is allocated or unused. Table 1 is pseudo-code for an exemplary data structures to store locally unique port **42** information. However, other data structures or layouts could also be used.

```
struct locally_unique_ports
{
        int port_number;
        flag status:1;   /* one bit flag, 0 = unused, 1 = allocated */
} gu_ports[MAX_GU];
int number_of_gu_ports;  /* number of locally unique ports allocated  */
```

Table 1.

The one or more locally unique ports **42** are allocated to protocols and applications in layered protocol stack **46** on the first network subdevice **6**. Upon receiving an unsuccessful assign response message **78**, the network subdevice may send another assign request message **76** for fewer ports. If the router **20** cannot allocate a large enough block of contiguous locally

McDonnell Boehnen Hulbert & Berghoff
300 South Wacker Drive, 32nd Floor
Chicago, IL 60606
(312) 913-0001

unique ports **42** for the first network subdevice **6**, it may send an assign response message **68** with a success code, but allocate fewer locally unique ports **42** than requested.

In an illustrative embodiment of the present invention, the router subsystem **20** allocates blocks of locally unique ports **42** to communications card **24**. However, other second network devices **7** could also be used to allocate locally unique ports **42** (e.g., a port server).

FIG. 14 is a block diagram illustrating a layout for a combination network address **112**. Combination network address layout **112** preferably includes a common external network address **44**, such as an IP **58** address, and a locally unique port **42** obtained by sending an assign request message **76** and receiving an assign response message **78** from a second network subdevice **7**. However, other layouts could also be used. The first network subdevices **6** use combination network address **112** for communications with external second network **14**. Common external network address **44** identifies the first private computer network **10** to an external second computer network **14**.

As is known in the art, to identify separate data streams, TCP **66** provides a source port field **114** and a source address field **116** in a TCP header **113**. Since local or default port identifiers are selected independently by each TCP **66** stack in a network, they are typically not unique. To provide for unique addresses within each TCP **66**, a local Internet address identifying TCP **66** can be concatenated with a local port identifier and a remote Internet address and a remote port identifier to create a "socket" that will be unique throughout all networks connected together. Sockets are known to those skilled in the networking arts.

In an illustrative embodiment of the present invention, the source port in a TCP header **113** is given a locally unique port **42** obtained with RSIP **64** and given a common external network address **44**. Together they uniquely identify applications and protocols on the first

McDonnell Boehnen Hulbert & Berghoff
300 South Wacker Drive, 32nd Floor
Chicago, IL 60606
(312) 913-0001

network subdevices **6** on first private computer network **10** to the second external computer network (e.g., **14** or **15**) with a value conceptually similar to the socket used by TCP **66**.

As is also known in the art, UDP **68** also has a source port field **118** in a UDP header **117**. The UDP source port field **118** is an optional field; when used, it indicates a port of the sending process, and may be assumed to be the port to which a reply should be addressed in the absence of any other information. If not used, a value of zero is inserted. A UDP header **117** also has a source address field **120**. A locally unique port can also be used in a UDP header **117**.

In an illustrative embodiment of the present invention, RSIP **70** is used to create combination network address **44** that is used in TCP header **113** and UDP header **117** fields. In another embodiment of the present invention, the combination network address **44** is stored in other message header fields understood by the router **20** (i.e., non-IP **58**, TCP **66**, or UDP **68** fields) and the second computer network **14**.

The router **20** maintains a port-to-internal network address table **122** as locally unique ports **42** are allocated. The router **20** also has an internal table **124** indicating internal network addresses for all network subdevices **5** on the first private computer network **10**. In an illustrative embodiment of the present invention, the internal network addresses for the first private computer network **10** are IP **58** addresses; however, other internal network addresses could also be used, such as a Medium Access Control ("MAC") protocol address. As an example, communications card **24** may have an internal IP address of 10.0.0.1, while the router **20** has an internal IP address of 10.0.0.2. The internal addresses are not published on the external computer network **14**.

FIG. 15 is a block diagram illustrating a port-to-internal address table **122** layout maintained by the router **20**. However, other layouts and more or fewer rows and columns could

alternatively be used . Port-to-internal address table **122** layout has three columns: an internal-network-address column **126**, a lowest-port column **128**, and a number-of-ports column **130**. However, more or fewer columns or other table layouts could also be used. The first row of the table **132** indicates that a first network subdevice **6** (e.g., a communications card **24**) has been allocated ports 1-32 for use with internal network address 10.0.0.1. A second network subdevice **7** (e.g., a the router **20**) uses ports 100-116 with internal network address 10.0.0.2. An internal network address may have several entries in port-to-internal address table **122**.

Realm Specific Internet Protocol

FIG. 16 is a flow diagram illustrating a method **134** for implementing RSIP in NAS **2**. At step **136**, a first network subdevice **6** on NAS **2** requests a common external address **44** and one or more locally unique ports **42** from a second network subdevice **7** on the NAS **2** with a first protocol **13**. The locally unique ports **42** are used in protocol layers in the layered protocol stack **42** on the first network subdevice **6**. In addition, the locally unique ports **42** are used to create a combination network address **112** comprising a locally unique port **42** and a common external address **44** to communicate with a second external computer network **14** without address translation. At step **138**, the first network subdevice **6** receives the common external address **44** and one or more locally unique ports **42** from the second network subdevice **7**. At step **140**, the first network subdevice **6** constructs one or more combination network addresses **112** using the one or more locally unique ports **42** and a common external network address **44** used to identify the NAS **2** to the second external computer network **14**.

In an illustrative embodiment of the present invention, the first network subdevice **6** is communications card **24**, the second network subdevice **7** is the router **20**, the first protocol **13** is

McDonnell Boehnen Hulbert & Berghoff
300 South Wacker Drive, 32nd Floor
Chicago, IL 60606
(312) 913-0001

RSIP **70**, the second external computer network **14** is the Internet or a private network. The combination network address includes a common IP **58** address (e.g., common network address **44**) identifying network subdevices **6** and **7** on NAS **2** to a second external computer network **14**. However, the present invention is not limited to the networks, network devices, network addresses or protocols described and others may also be used.

The ports **42** are used for entities such as protocols and applications in layered protocol stack **42** on network device and are locally unique on NAS **2**. The locally unique ports **42** will identify a network subdevice on NAS **2**. After allocation with method **130**, a network subdevice uses a locally unique port **42** in a protocol layer in layered protocol stack **42**. As is illustrated in FIG. 15, first network subdevice **6** with internal IP **58** address 10.0.0.1 is assigned thirty-two locally unique ports in the range of 1-32. The first network subdevice **6** may assign locally unique port-2 to TCP **66** to use as a source port. The combination network address **112** illustrated in FIG. 14 is then assigned to TCP **66** on the first network subdevice **6** for communications with an external network (e.g., **14** or **15**). Other locally unique ports **42** are assigned to other protocols and applications in layered protocol stack **42** on a network subdevice **6**.

In one embodiment of the present invention, locally unique ports are assigned to protocol layers in layered protocol stack **42** when a network device boots. In another embodiment of the present invention, locally unique ports are assigned to protocol layers in layered protocol stack when a protocol layer makes a request for an external network (e.g., **14** or **15**). In yet another embodiment of the present invention, locally unique ports are assigned dynamically or on-the-fly in an individual protocol layer as a protocol layer makes a request for an external network (e.g., **14** or **15**).

McDonnell Boehnen Hulbert & Berghoff
300 South Wacker Drive, 32nd Floor
Chicago, IL 60606
(312) 913-0001

The locally unique ports **42** with common external network address **44**, together forming combination network address **112**, uniquely identify a network subdevice **6** to an external network **14** without translation.

FIG. 17 is a flow diagram illustrating a method **140** for implementing RSIP. At step **142**, a communication is sent from a first network subdevice **6** on a NAS **2** to a second network subdevice **7** on the NAS **2**. The communication is for a second external network **14** and includes a combination network address **112** identifying the first network subdevice **6** on the NAS **2**. The combination network address **112** is constructed with method **130** (FIG. 16) and includes a locally unique port **42** and a common external address **44** to identify the NAS **2** to the second external network **14**. At step **144**, the second network subdevice **7** routes the request from the NAS **2** to the second external network **14**. At step **146**, the second network subdevice **7** on the NAS **2** receives a response communication from the external second computer network **14** at the external network address **44** identifying the NAS **2** from the combination network address **112**. At step 148, the second network subdevice **7** on the NAS **2** routes the response communication to the first network subdevice **6** on the NAS **2** using the locally unique port **42** from the combination network address **112**.

In an illustrative embodiment of the present invention, the first network subdevice **6** is a communications card **24**, the second network subdevice is a the router **20**, the NAS **2** is a stub network, and the second computer network is the Internet or a private network. The combination network address **112** includes a locally unique port **42** obtained with RSIP **70** and an external IP **58** address **44** for an external network **14** such as the Internet, an intranet, or another computer network. However, the present invention is not limited to the networks, network devices, network address or protocol described and others may also be used.

McDonnell Boehnen Hulbert & Berghoff
300 South Wacker Drive, 32nd Floor
Chicago, IL 60606
(312) 913-0001

The method **140** (FIG. 17) is illustrated with a specific example using TCP **66**/IP **58** layers from layered protocol stack **42**. However, other protocol layers in layered protocol stack **42** could also be used. At step **142**, the first network subdevice **6** sends a TCP **66** communication to the router **20**, for example, a TCP **66** communication for the router **20** at external IP **58** address 192.200.20.3 on second computer network **30**. Table 2 illustrates an example of a communication data packet sent at step 142.

IP **58** Header	TCP **66** Header
SRC IP: 198.10.20.30	SRC Port: 2
DST IP: 192.200.20.3	DST Port: 80

Table 2.

The source IP **58** address is common external network address **44** (e.g., 198.10.20.30) and the source port is locally unique port 2 obtained via RSIP **70** with the method **130** and assigned to TCP **66**. In one embodiment of the present invention, the locally unique port 2 for TCP **66** is requested and assigned when the first network subdevice **6** is booted. In another embodiment of the present invention, the locally unique port 2 is assigned when a protocol layer in the layered protocol stack initiates the communication with the external network **14**. The locally unique port along with the common external address **44** comprise the combination network address **112**. The destination IP address is 192.200.20.3 for the router **20** (FIG. 2) on the second external network **30** and the destination port is well known Internet port **80**. When the communication reaches the link layer **48**, in the layered protocol stack **42**, an outer IP **58** header is added to route the communication to the router **20**. The local internal network address (e.g., 10.0.0.x) for a network subdevice for internal communications is maintained in the link layer **48**. Table 3 illustrates an exemplary data packet with an outer IP **58** header added for the router **20**.

Outer IP **58** header	Inner IP **58** header	TCP **66** header
SRC IP: 10.0.0.1	SRC IP: 198.10.20.30	SRC Port: 2
DST IP: 10.0.0.7	DST IP: 192.200.20.3	SRC Port: 80

Table 3.

The link layer **48** adds the outer IP **58** header including a source IP **58** address for the first network subdevice **6** of 10.0.0.1 and a destination IP **58** address of 10.0.0.7 for the router **20**. At step **144**, the router **20** receives the communication data packet, strips the outer IP **58** header, and sends the communication data packet to the external network **14**.

At step **146**, the router **20** receives a response communication packet from an external network (e.g., **30**). An example of a response data packet is illustrated in Table 4.

IP **58** Header	TCP **66** Header
SRC IP: 192.200.20.3	SRC Port: 80
DST IP: 198.10.20.30	DST Port: 2

Table 4.

The router **20** receives the response packet from the external second network **14** at step **146** with the destination IP **58** address, common external network address 198.10.20.30 and the destination port set to locally unique port 2. The router **20** uses port-to-internal network address table (FIG. 15) to map destination port 2 to the internal IP **58** address 10.0.0.1 for first network device **6**. The router **20** adds an outer IP **58** header to route the response data packet back to the first network subdevice **6**. Table 5 illustrates an exemplary response packet with outer IP **58** header added by the router **20**.

Outer IP **58** header	Inner IP **58** header	TCP **66** header
SRC IP: 10.0.0.7	SRC IP: 192.200.20.3	SRC Port: 80
DST IP: 10.0.0.1	DST IP: 198.10.20.30	SRC Port: 2

Table 5.

The outer IP **58** header has a source internal IP **58** address of 10.0.0.7 for the router **20** and a destination internal IP **58** address of 10.0.0.1 for the first network subdevice **6** on the computer network **10**. At step **148**, the router **20** routes the response data packet to the first

network subdevice **6** with the outer IP **58** header. The link layer **48** in the layered protocol stack **42** strips the outer IP **58** header and forwards the response data packet to the network layer **50**.

The first network subdevice **6** sends a communication to an external network **14** and receives a response communication from the external network **14** using RSIP **70** and the locally unique port **42** allocated with RSIP **70**. The router **20** does not translate any source/destination IP **58** addresses or source/destination ports. Thus, RSIP is accomplished without network address translation at the router **20**.

An illustrative embodiment of the present invention is described with respect to a single common external network address **44** identifying multiple network subdevices **6** and **7** on NAS **2** and used in the combination network address **112** with a locally unique port **42**. However, the present invention is not limited to a single common external network address and can also be practiced with a multiple common external network addresses.

RSIP using method **134** (FIG. 16) and method **140** (FIG. 17) removes the computational burden of NAT at the router **20** and allows multiple network subdevices **6** to use a single or a small number of external network addresses **44** known to an external network **14** such as the Internet or an intranet. Instead of providing NAT, the router **20** routes data packets from a first network subdevice **6** on NAS **2** to a second external computer network **14** using the combination network address **112**. In addition, the router **20** is no longer required to support multiple application protocols from the layered protocol stack **42**.

The router **20** also routes data packets from the second external computer network **14** back to a first network subdevice **6** on the NAS **2** using the locally unique port **42** in the combination network address **112**. The router **20** is no longer required to replace an internal network address **10** with an external network address **44** for outbound traffic, nor to replace an

McDonnell Boehnen Hulbert & Berghoff
300 South Wacker Drive, 32nd Floor
Chicago, IL 60606
(312) 913-0001

external network address **44** with an internal network address **11** for inbound traffic. Thus, RSIP of the present invention removes the computational burden of NAT from the router **20** and does not violate the Internet principal of providing end-to-end transmission of data packets between network devices without alterations. This allows end to end protocols, such as IPsec, to work between the NAS **2** and the external network **14**.

An embodiment of the architecture of the present invention is an IP telephony system. In this case, NAS **2** is an IP Telephony Gateway system, with the communications cards **24** acting as media translators between the public-switched telephone network ("PSTN"; not shown in Figure 2) and the internal IP network **10**. Data application **32** provides media translation (gateway) functionality, while a device control application **30** allows each media device (per card) to be controlled remotely. More specifically, each communications card **24** appears as a MEGACO-compliant (*See* Cuervo et al., "*Megaco Protocol*," Internet Draft <draft-ietf-megaco-protocol-06.txt, Feb. 8, 2000, incorporated herein by reference) Media Gateway ("MG"), and the IP device control element **31** on the external IP signaling network **29** is a MEGACO-compliant Media Gateway Controller ("MGC"). The signaling/control communications between the MGC and MG in MEGACO is indicated in Figure 2 by the dashed line between device control application **30** on communications card **24**, and the IP control device **31** on the external signaling network **29**.

Similarly, data application **32** on each card **24** provides the media capability of the MG. The specific applications for this example includes real-time transport protocol ("RTP") for transport of the media on the IP network. For more information on RTP, see H. Schulzrinne, et al., *RTP: A Transport Protocol for Real-Time Applications*, RFC-1889, incorporated herein by reference. This application communicates with a peer application on an external media device

McDonnell Boehnen Hulbert & Berghoff
300 South Wacker Drive, 32nd Floor
Chicago, IL 60606
(312) 913-0001

33. The media communications between the MG and its IP peer are indicated in Figure 2 by the dashed line between data application **32** on communications card **24**, and IP media device **33** on external IP data network **27**.

In the case of both MEGAGO and RTP, external communications with the internal MG uses IP. Therefore the internal MG must be IP-addressable. As shown in the example configuration, each communications card **24** has an IP interface **26** to support IP communications with other IP devices, such as the MGC or external media device. However, the IP interface **26** on each card **24** provides only an internal (private) IP address. The external (public) IP interfaces are provided by router subsystem **20**. Therefore, the collection of communications cards **24** (distributed among multiple chassis 18 in this example) comprises a stub network. By implementing an RSIP gateway **38** on router subsystem **20**, and an RSIP host **28** on each of the communications cards **24**, external IP devices, such as and MGC or external MG (peer media device), can communicate directly with the internal, card-based MG over one or more external IP networks connected at the router subsystem **20**.

There are a number of alternative configurations for the implementation of RSIP in a NAS such as the one illustrated in Figure 2. If the router subsystem **20** provides multiple external IP interfaces **12**, then a single RSIP gateway **38** must be able to distinguish among them in order to properly route packets across them. Alternatively, a separate RSIP gateway **38** may be implemented on each external IP interface **12**. In another preferred embodiment, the functionality of the RSIP gateway **38** is decomposed in such a way as to provide a single, common management component, and separate mapping components (one for each external IP interface **12**). Similarly, each communications card **24** may implement one or multiple RSIP hosts **28**, where the choice may depend upon the number of IP interfaces **12** on each card **24**.

McDonnell Boehnen Hulbert & Berghoff
300 South Wacker Drive, 32nd Floor
Chicago, IL 60606
(312) 913-0001

Finally, the RSIP gateway **38** may reside on a subsystem other than the router subsystem **20**. The only requirement is that RSIP gateway **38** resides between the internal and external IP interface(s).

Address mapping/sharing may also be required in cases where two different address spaces must be bridged. The embodiments presented herein assume that the address spaces are an internal, private IP network and external, public IP networks. The method of using RSIP **70** in a NAS **2**, however, applies equally well for bridging networks using different versions of IP **58**. For example, the invention could bridge an internal IPv4 network and an external IPv6 network, an internal (private) IPv4 network and an external (public) IPv6 network, an internal (private) IPv6 network and an external (public) IPv4 network, or an internal (private) IPv6 network and an external (public) IPv6 network.

The various embodiments of the present invention described above offer several advantages over the prior art. Network address translation and the large computational burden is removed from a router and distributed to individual network devices using a port allocation protocol to allocate locally unique ports and globally unique addresses. RSIP with port translation does not violate the Internet principal that recommends that packets flow end-to-end between network devices without changing the contents of any packet along a transmission route. Illustrative embodiments of the present invention can support multi-casting with a router serving as a proxy for internal network devices that wish to join an existing multicast session. Illustrative embodiments of the present invention can also be used to support Virtual Private Networks ("VPNs").

RSIP also allows a local network to efficiently switch between external network service providers (e.g., Internet service providers) by changing the common external address for an

external network assigned to a local network. RSIP also allows a local network to purchase a smaller block of external network addresses, providing a cost savings on the local network.

The various embodiments of the present invention described above offer several advantages over the prior art. Network address translation and the large computational burden is removed from a router and distributed to individual network devices using a port allocation protocol to allocate locally unique ports. RSIP does not violate the Internet principal that recommends that packets flow end-to-end between network devices without changing the contents of any packet along a transmission route, which breaks some protocols. Moreover, for a system such as the one illustrated in Figure 2, the ability to terminate an IPsec connection (security association) at the device control application and/or data application on any communications card might be a requirement. RSIP introduces no hindrance to such a required capability, and was, in fact, developed specifically to allow for it. If NAT were used instead to provide the address mapping/sharing functionality for such a system, IPsec connections to the individual IP-addressable sub-components of the system would be impossible. RSIP is the only method for providing the simultaneous capabilities of address mapping/sharing and end-to-end IPsec. This applies equally for any application or protocol that requires end-to-end connectivity (i.e., strict disallowance of packet modification by intermediate routers, forwarders, etc.). The methods of the present invention are useful with IPsec as previously described in U.S. Application No. 09/270,967, filed March 17, 1999, incorporated herein by reference.

It should be understood that the programs, processes, methods and apparatus described herein are not related or limited to any particular type of computer or network apparatus (hardware or software), unless indicated otherwise. Various types of general purpose or

specialized computer apparatus may be used with or perform operations in accordance with the teachings described herein.

In view of the wide variety of embodiments to which the principles of the present invention can be applied, it should be understood that the illustrated embodiments are exemplary only, and should not be taken as limiting the scope of the present invention. For example, the steps of the flow diagrams may be taken in sequences other than those described, and more or fewer elements may be used in the block diagrams.

The claims should not be read as limited to the described order or elements unless stated to that effect. In addition, use of the term "means" in any claim is intended to invoke 35 U.S.C. §112, paragraph 6, and any claim without the word "means" is not so intended. Therefore, all embodiments that come within the scope and spirit of the following claims and equivalents thereto are claimed as the invention.

WHAT WE CLAIM IS:

1. A method of implementing Realm Specific Internet Protocol in a network access system comprising a plurality of network subdevices connected by a network, the method comprising the steps of:

(a) requesting by a first network subdevice using a first protocol, a common external network address and one or more ports from a second network subdevice to identify a first network subdevice during communications with an external computer network;

(b) receiving the common external network address and an identifier of the one or more ports at the first network subdevice from the second network subdevice ;

(c) updating entries in an address-to-address table maintained by the second network device to reflect assignment of the common external network address and one or more ports to the first network subdevice; and

(d) creating a combination network address for the first network subdevice with the identifier of the one or more ports and the common external network address, the combination network address identifying the first network subdevice for communications with the external computer network..

2. A computer readable medium having stored therein instructions for causing a central processing unit to execute the method of claim 1.

3. The method of claim 1 further comprising:

(a) sending a request from the first network subdevice to the second network subdevice;

McDonnell Boehnen Hulbert & Berghoff
300 South Wacker Drive, 32nd Floor
Chicago, IL 60606
(312) 913-0001

(b) routing the request from the second network subdevice to the external computer network;

(c) receiving a reply at the second network subdevice on the common external network address for the network access system; and

(d) routing the reply from the second network subdevice to the first network subdevice using the locally unique port from the combination network address.

4. The method of claim 1 wherein the first protocol is a Realm Specific Internet Protocol comprising a Realm Specific Internet Protocol assign request message, a Realm Specific Internet Protocol assign response message, and a combination network address involving a locally unique port and a common external network address.

5. The method of claim 1 wherein the common external network address is an Internet protocol address.

6. The method of claim 1 wherein the first network subdevice is a communications card.

7. The method of claim 6 wherein the communications card comprises a Realm Specific Internet Protocol host and an Internet protocol interface.

8. The method of claim 7 wherein the communications card further comprises a data application and a device control application.

9. The method of claim 1 wherein the second network subdevice is a router or a port server.

10. The method of claim 1 wherein the second network subdevice comprises a Realm Specific Internet Protocol gateway and a plurality of Internet protocol interfaces.

11. The method of claim 1 wherein the external computer network is any of the Internet, an intranet or a public-switched telephone network.

12. The method of claim 1 wherein the common external network address is an Internet protocol address.

13. The method of claim 1 wherein the plurality of subdevices on the network access system comprise a local area network and the external network is any of the Internet or an intranet.

14. A network access device, comprising in combination:

(a) a first network;

(b) a first network subdevice comprising a network clienton the first network , wherein the first network subdevice has a first network address for communicating with other network subdevices and requests from a second network subdevice allocation of a second

network address and one or more ports for communicating with a plurality of network devices on a second network; and

(c) a second network subdevice on the first network comprising a network address server for allocating a second network address and one or more ports to the first network subdevice, wherein the second network subdevice has a first network address for communicating with other network subdevices on the first network and a second network address for communicating with a plurality of network devices on a second network, and wherein the network address server is used to allocate the second network address to the first network subdevice on the first network.

15. The network access device of claim 14 wherein the first network is a private Internet Protocol network.

16. The network access device of claim 14 wherein the second network is a public network.

17. The network access device of claim 14 wherein the first network address of the first network subdevice is a private network address

18. The network access device of claim 14 wherein the first network address of the second network subdevice is a private network address and the second network address of the second network subdevice is a public network address.

McDonnell Boehnen Hulbert & Berghoff
300 South Wacker Drive, 32nd Floor
Chicago, IL 60606
(312) 913-0001

19. The network access device of claim 14 wherein the first network subdevice further comprises an IP interface and the client of the first network subdevice is a Realm Specific Internet Protocol host.

20. The network access device of claim 14 wherein the second network subdevice further comprises an IP interface and the network address server of the second network subdevice is a Realm Specific Internet Protocol gateway.

21. The network access device of claim 14 wherein the first network subdevice is a communications card.

22. The network access device of claim 21 wherein the communications card is a modem card.

23. The network access device of claim 14 wherein the first network subdevice further comprises a data application and a device control application.

24. The network access device of claim 17 wherein the private network address of the second network subdevice is an Internet protocol address.

25. The network access device of claim 14 wherein the second network subdevice is a router subsystem.

McDonnell Boehnen Hulbert & Berghoff
300 South Wacker Drive, 32nd Floor
Chicago, IL 60606
(312) 913-0001

26. The network access device of claim 17 wherein the public network address of the second network subdevice is an Internet protocol address.

27. The network access device of claim 16 wherein the second network is any of the Internet or an intranet.

28. The network access device of claim 23 wherein the network access device is an Internet telephony gateway system.

29. The network access device of claim 28 wherein the data application provides media translation functionality and the device control application provides for remote control of the first network subdevice by a network device on the second network.

30. The network access device of claim 28 wherein the first network subdevice is a MEGACO-compliant media gateway.

31. The network access device of claim 28 wherein the second network comprises an external Internet Protocol signaling network having an Internet Protocol control device and an external Internet Protocol data network having an Internet Protocol media device.

32. The network access device of claim 31 wherein the Internet Protocol control device on the external Internet Protocol signaling network is a MEGACO-compliant media gateway controller.

33. The network access device of claim 14 wherein the first private network subdevice and the second private network subdevice are cards in a rack having a common backplane.

ABSTRACT

A method to support the assignment of a globally unique network public address and, optionally, a number of locally unique ports to a first private network subdevice having a private network address on a network access system from a second private network subdevice having a public network address using Realm Specific Internet Protocol, wherein the public network address is used by the first private network device to communicate with network devices on an external network.

McDonnell Boehnen Hulbert & Berghoff
300 South Wacker Drive, 32nd Floor
Chicago, IL 60606
(312) 913-0001

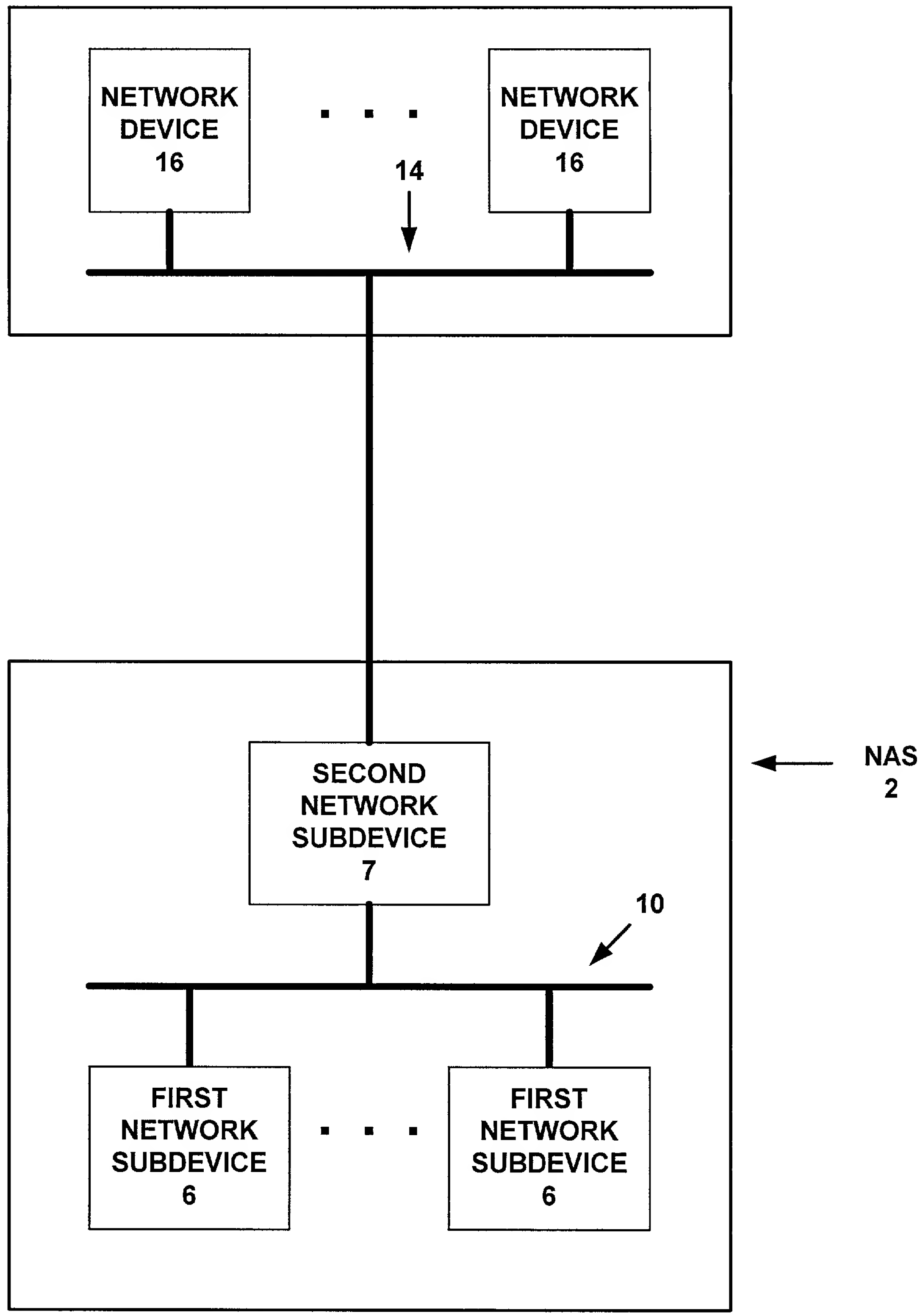
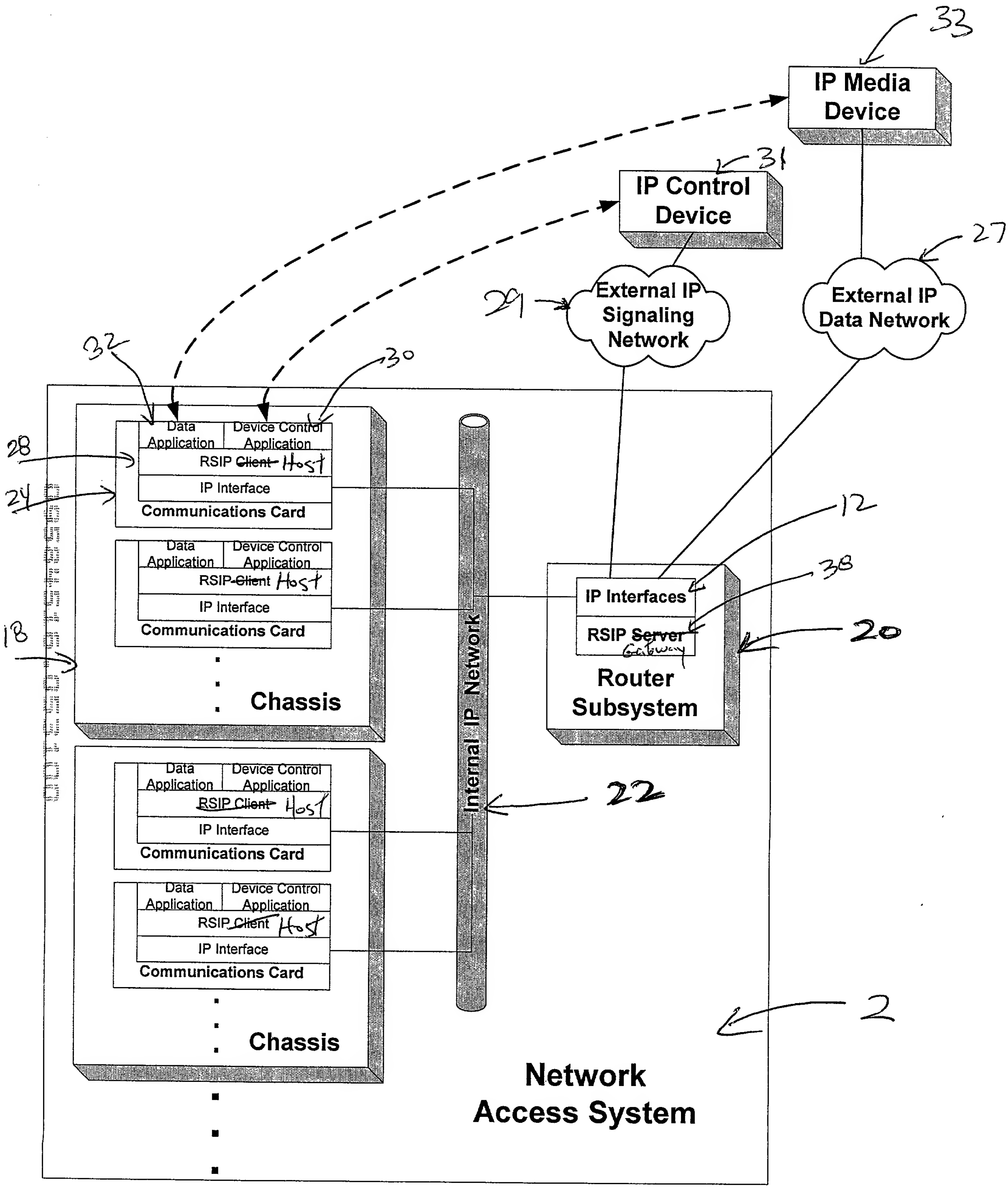

NETWORK DEVICE 16
NETWORK DEVICE 16
14
SECOND NETWORK SUBDEVICE 7
10
NAS 2
FIRST NETWORK SUBDEVICE 6
FIRST NETWORK SUBDEVICE 6

FIGURE 1



33
IP Media Device
31
IP Control Device
27
29
External IP Signaling Network
External IP Data Network
32
30
Data Application
Device Control Application
28
RSIP ~~Client~~ Host
24
IP Interface
Communications Card
Data Application
Device Control Application
RSIP ~~Client~~ Host
IP Interface
Communications Card
18
Chassis
12
38
IP Interfaces
RSIP ~~Server~~ Gateway
20
Router Subsystem
Internal IP Network
22
Data Application
Device Control Application
RSIP ~~Client~~ Host
IP Interface
Communications Card
Data Application
Device Control Application
RSIP ~~Client~~ Host
IP Interface
Communications Card
Chassis
2
Network Access System

FIGURE 2

FIG. 3

RSIP PROTOCOL STACK

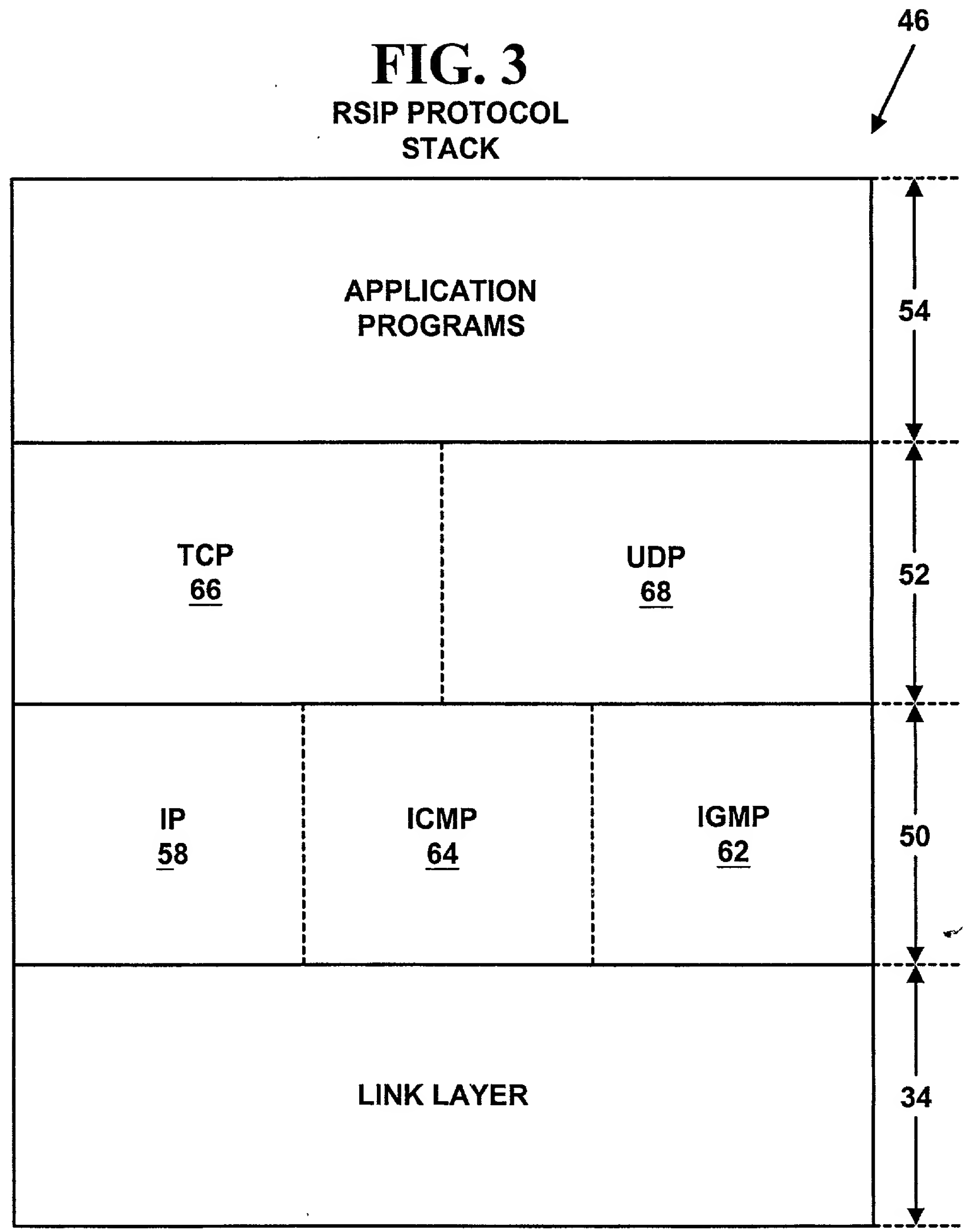


FIG. 4

REALM SPECIFIC INTERNET PROTOCOL (RSIP) MESSAGES

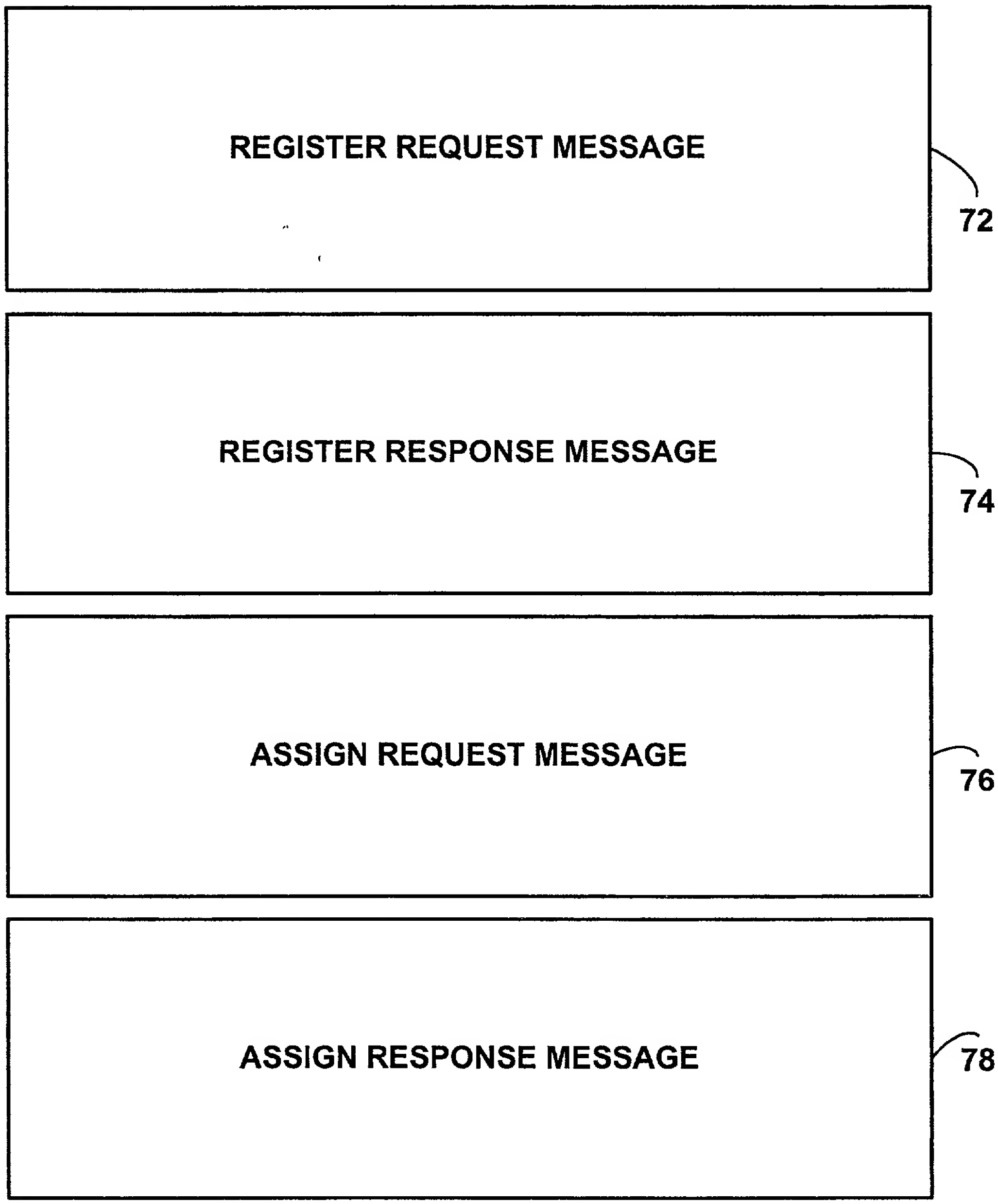


FIG. 5

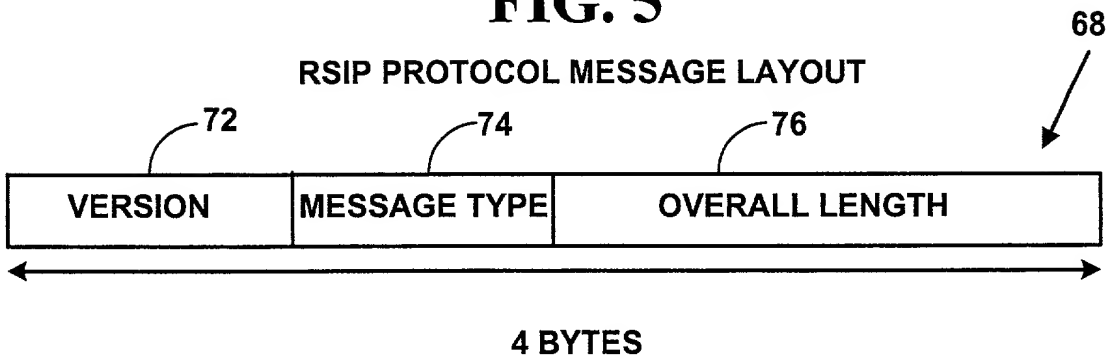


FIG. 6

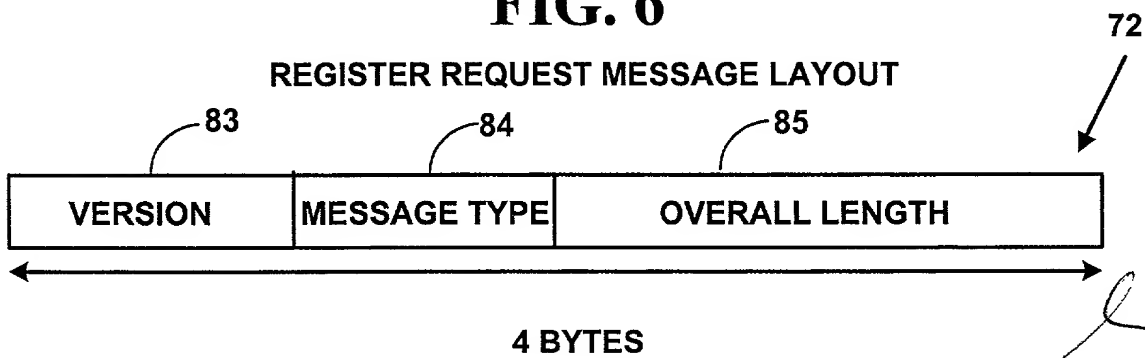


FIG. 7

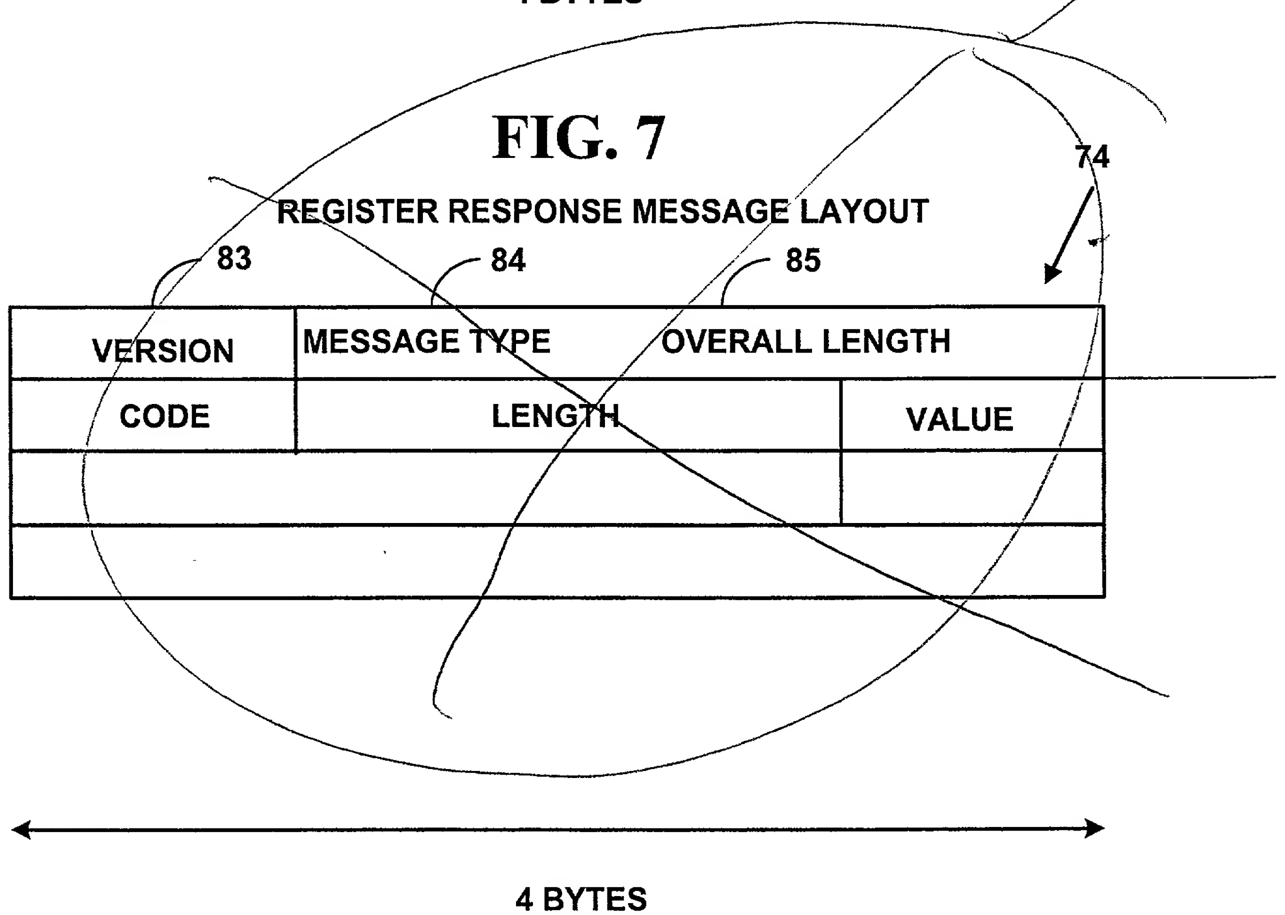


FIGURE 7
REGISTER RESPONSE MESSAGE LAYOUT

VERSION	MESSAGE TYPE	OVERALL LENGTH		
CODE	LENGTH		VALUE	CLIENT ID
CLIENT ID VALUE (CONT'D)			CODE	FLOW
LENGTH		LOCAL	REMOTE	POLICY
CODE	LENGTH	VALUE		RSIP METHOD
CODE	LENGTH	VALUE		TUNNEL TYPE

4 BYTES

FIGURE 8
ASSIGN REQUEST MESSAGE LAYOUT

VERSION	TYPE	LENGTH		
CODE	LENGTH		VALUE	CLIENT ID
CLIENT ID VALUE (CONT'D)			CODE	LOCAL ADDRESS
LENGTH		TYPE		
VARIABLE LENGTH VALUE				
METHOD DEPENDENT FIELDS				

4 BYTES

FIGURE 9
RSA-IP ASSIGN REQUEST MESSAGE LAYOUT

CODE	LENGTH		TYPE	REMOTE
VARIABLE LENGTH VALUE				ADDRESS
CODE	LENGTH		NUMBER	REMOTE
VARIABLE NUMBER OF 2 BYTE PORT FIELDS				PORTS
CODE	LENGTH		VALUE	[LEASE TIME]
LEASE TIME VALUE CONT'D			CODE	[TUNNEL
LENGTH		VALUE		TYPE]

4 BYTES

FIGURE 10
RSAP-IP ASSIGN REQUEST MESSAGE LAYOUT

CODE	LENGTH		NUMBER	LOCAL
VARIABLE NUMBER OF 2 BYTE PORT FIELDS				PORTS
CODE	LENGTH		TYPE	REMOTE
VARIABLE LENGTH VALUE				ADDRESS
CODE	LENGTH		NUMBER	REMOTE
VARIABLE NUMBER OF 2 BYTE PORT FIELDS				PORTS
CODE	LENGTH		VALUE	[LEASE TIME]
LEASE TIME VALUE CONT'D			CODE	[TUNNEL
LENGTH		VALUE		TYPE]

4 BYTES

FIGURE 11
ASSIGN RESPONSE MESSAGE LAYOUT

VERSION	TYPE	LENGTH		
CODE	LENGTH		VALUE	CLIENT ID
CLIENT ID VALUE (CONT'D)			CODE	
LENGTH		VALUE		BIND ID
BIND ID VALUE CONT'D		CODE	LENGTH	LOCAL
LENGTH CONT'D	TYPE	VARIABLE LENGTH VALUE		ADDRESS
METHOD DEPENDENT FIELDS				

4 BYTES

FIGURE 12
RSA-IP ASSIGN RESPONSE MESSAGE LAYOUT

CODE	LENGTH		TYPE	REMOTE
VARIABLE LENGTH VALUE				ADDRESS
CODE	LENGTH		NUMBER	REMOTE
VARIABLE NUMBER OF 2 BYTE PORT FIELDS				PORTS
CODE	LENGTH		VALUE	LEASE TIME
LEASE TIME VALUE CONT'D			CODE	TUNNEL
LENGTH		VALUE		TYPE

4 BYTES

FIGURE 13
RSAP-IP ASSIGN REQUEST MESSAGE LAYOUT

CODE	LENGTH	NUMBER	LOCAL
VARIABLE NUMBER OF 2 BYTE PORT FIELDS			PORTS
CODE	LENGTH	TYPE	REMOTE
VARIABLE LENGTH VALUE			ADDRESS
CODE	LENGTH	NUMBER	REMOTE
VARIABLE NUMBER OF 2 BYTE PORT FIELDS			PORTS
CODE	LENGTH	VALUE	LEASE TIME
LEASE TIME VALUE CONT'D		CODE	TUNNEL
LENGTH	VALUE		TYPE

4 BYTES

FIG. 14

112

44 COMBINATION NETWORK ADDRESS 42

EXTERNAL NETWORK ADDRESS (E.G., EXTERNAL IP ADDRESS)	LOCALLY UNIQUE PORT
198.10.20.30	1032

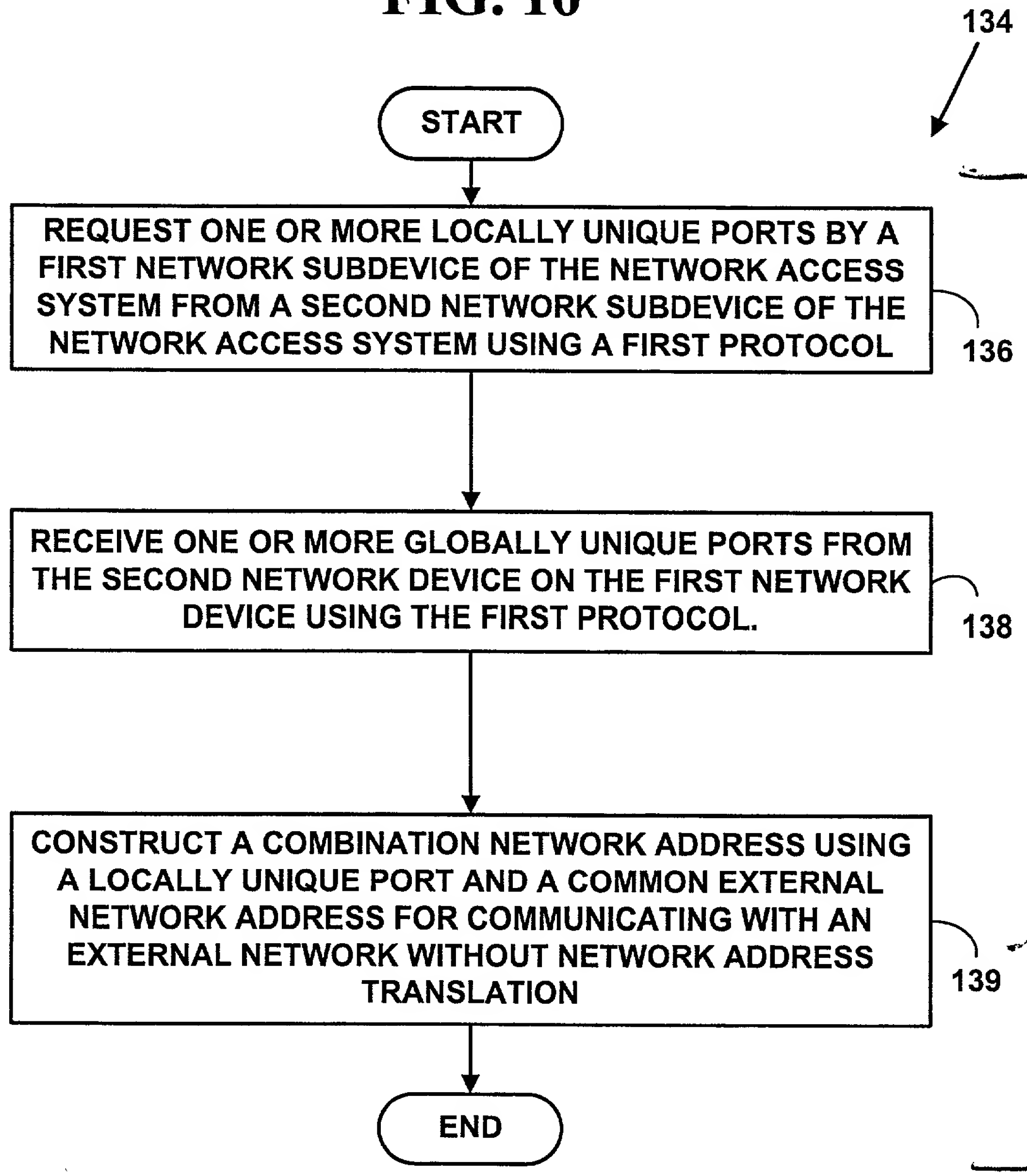
FIG. 15

122

126	128	130
INTERNAL NETWORK ADDRESS	LOWEST PORT	NUMBER OF PORTS
10.0.0.1	1026	32
10.0.0.3	1057	16

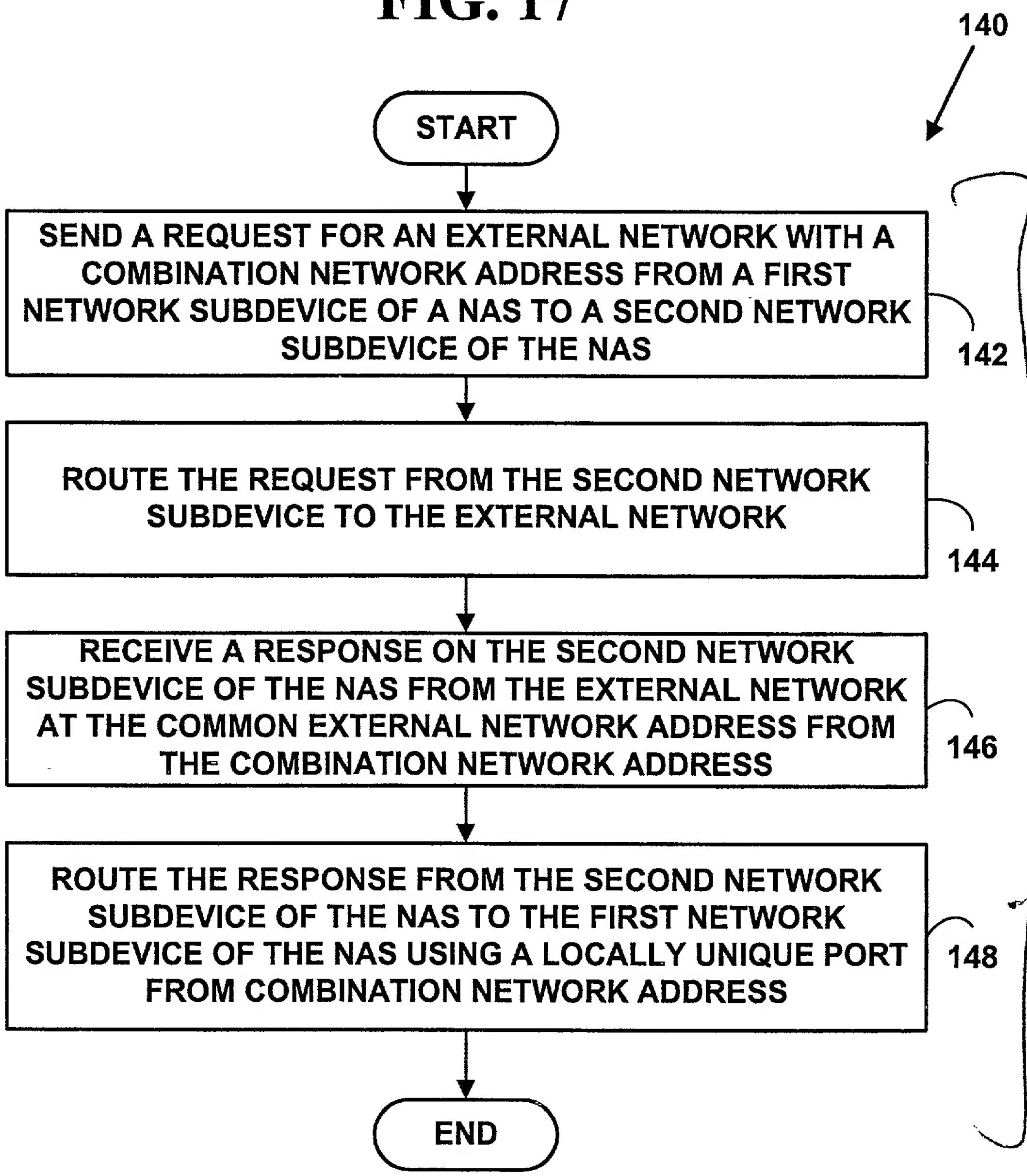
PORT-TO-INTERNAL-NETWORK
ADDRESS TABLE

FIG. 16



These will be rewritten for RAS rather than NAS

FIG. 17



These will be rewritten for RAS rather than NAS